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Integration of XML data in P2P E-commerce applications

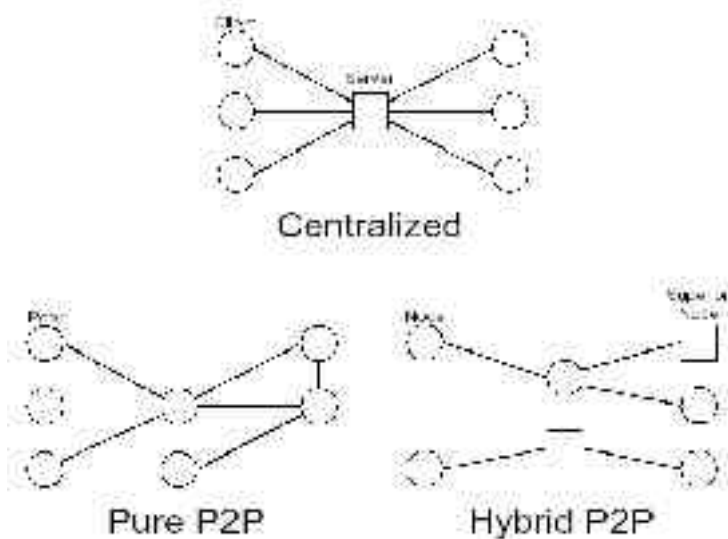
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Overview

1. Motivation – XML documents in e-commerce and needs of their integration
2. Approaches to data integration: materializing (*data transformation*) vs virtual data integration (*query reformulation*)
3. XML schema mapping – a language XDMMap
4. Automating mapping specification and query reformulation
5. Conclusion

Motivation – P2P E-commerce

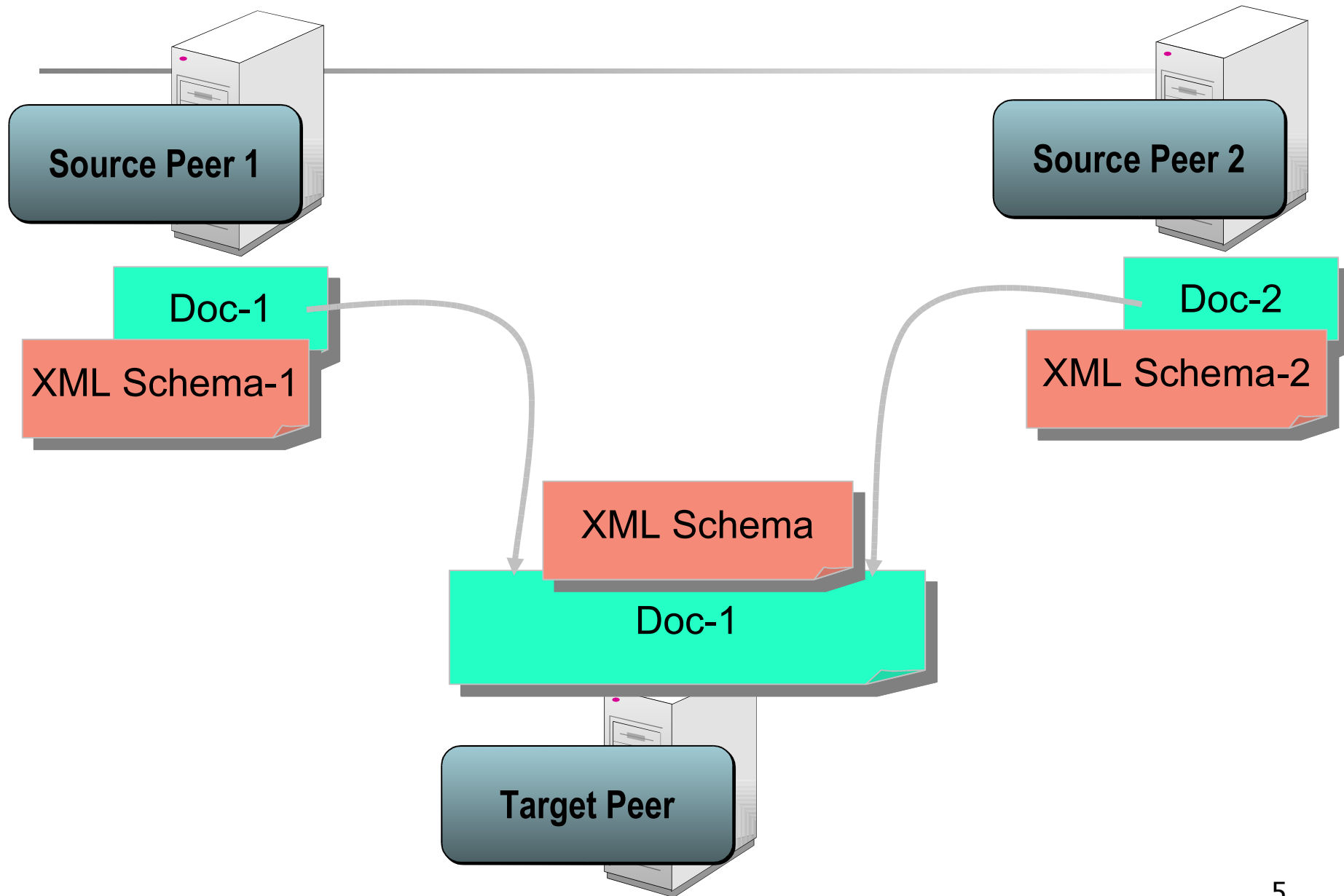


1. E-commerce applications need new solutions in the field of integration, interchange, and transformation of data across the global marketplace.
2. Traditional client-server systems are replaced by P2P systems that offer more flexibility.
3. Advanced data management features are expected to function automatically or semi-automatically.
4. A basic functionality of a system supporting e-commerce applications is the **integration** of external data sources or data services.

Sample scenario

- a newly hired employee requires a laptop for his workplace,
- the employee defines characteristics of the laptop and some additional conditions using an online website which is actually composed of several distinct and interacting Web Services,
- based on the specified order, catalog Web Service contacts Web Services of some hardware vendors and collects a list of offers,
- a price comparisons can be done and some additional negotiations can be performed (e.g. a service contract, prices, payment, etc.).

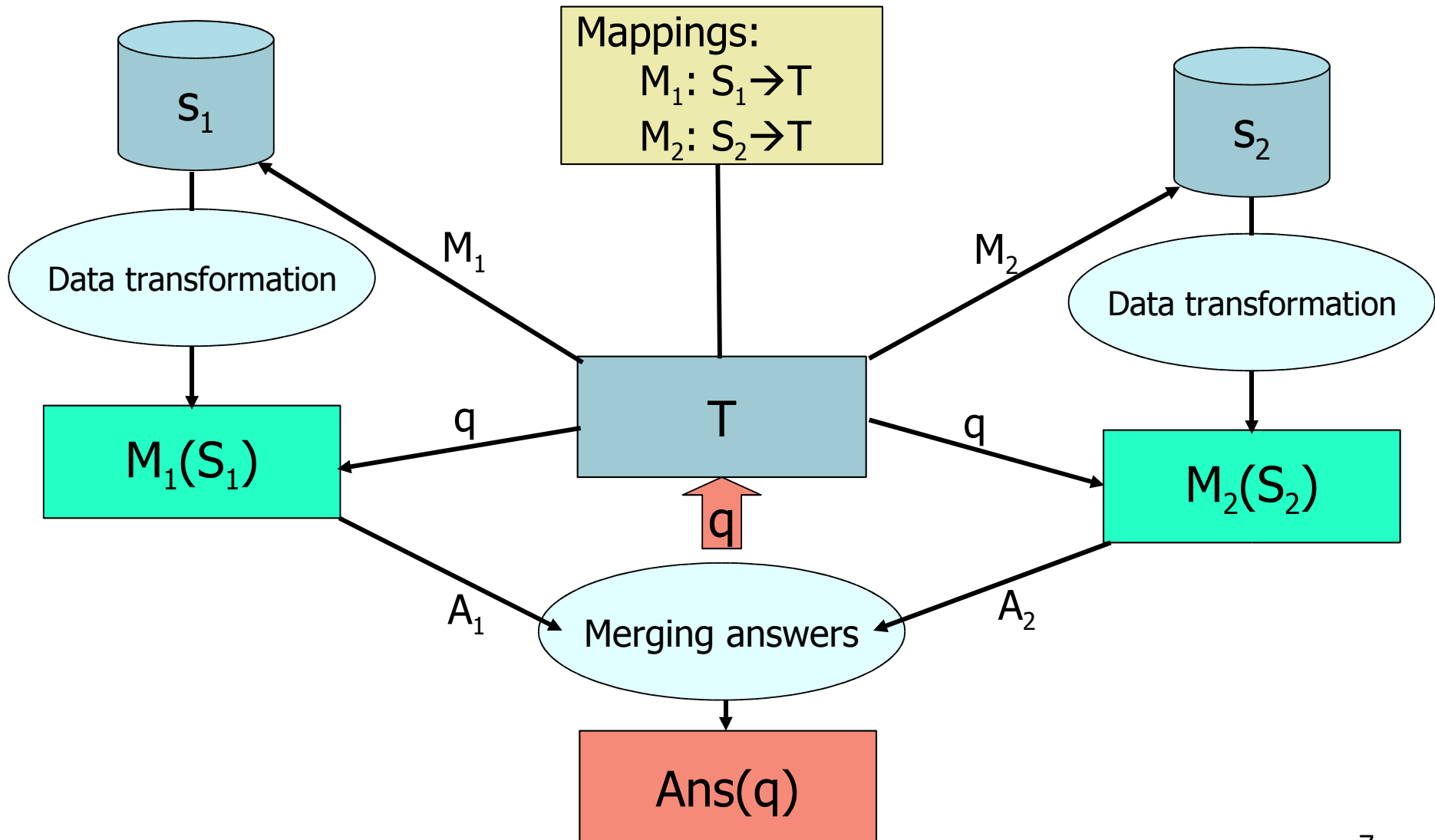
Integrating data from XML documents



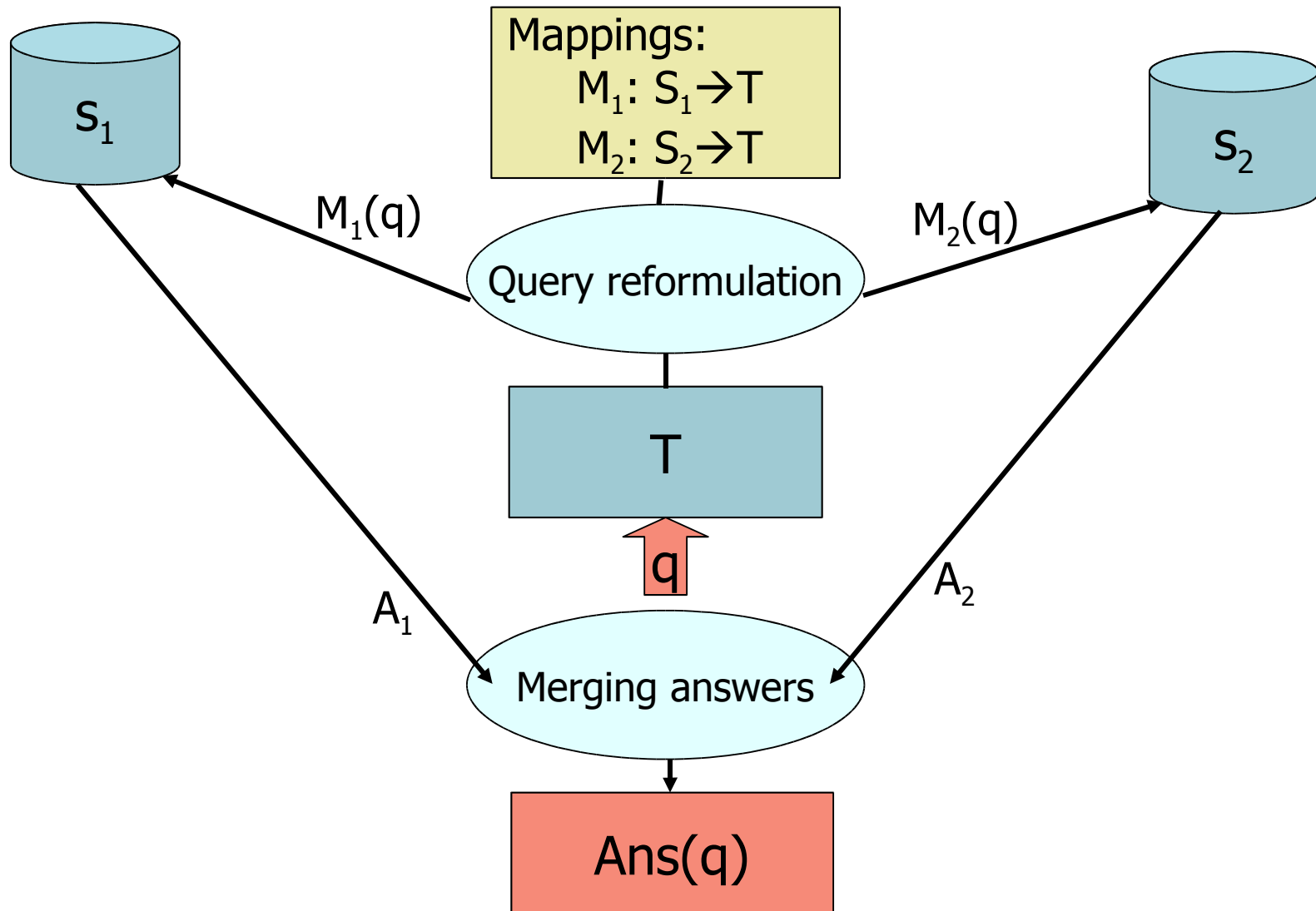
Data integration

1. System architecture
 - Client-server
 - P2P Data Management Systems (PDMS)
2. Strategy of integration
 - Data transformation (materializing integration)
 - Query reformulation (virtual integration)

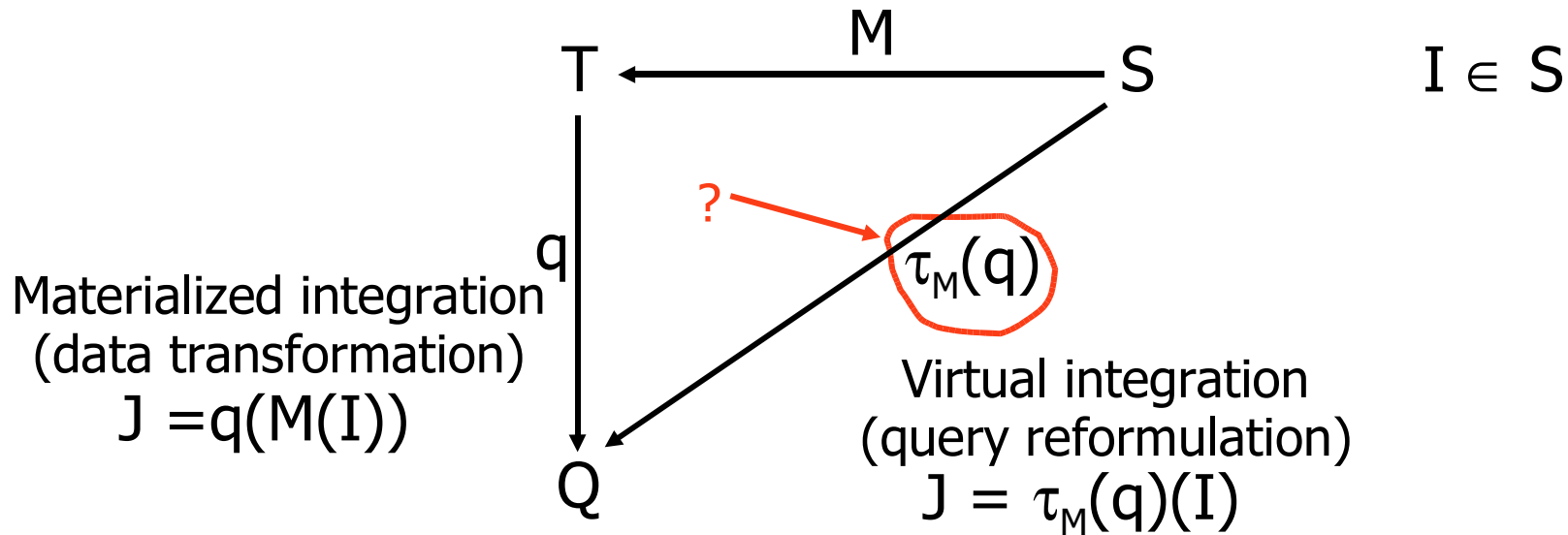
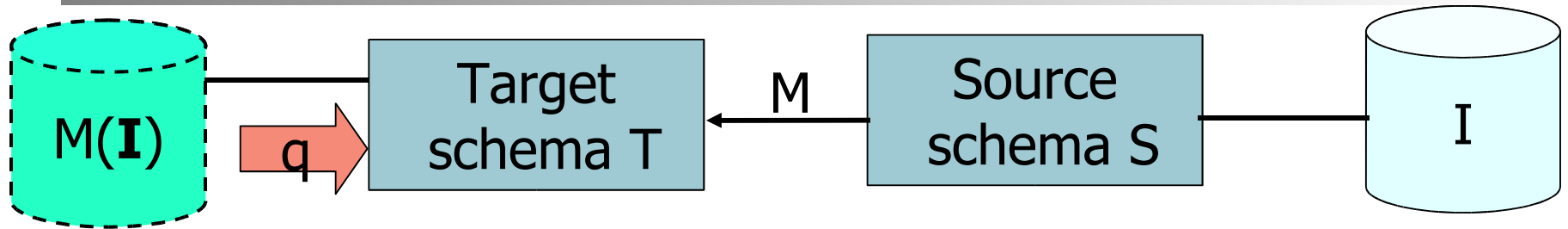
Data integration by data transformation



Data integration by query reformulation



Materialized vs Virtual Data integration

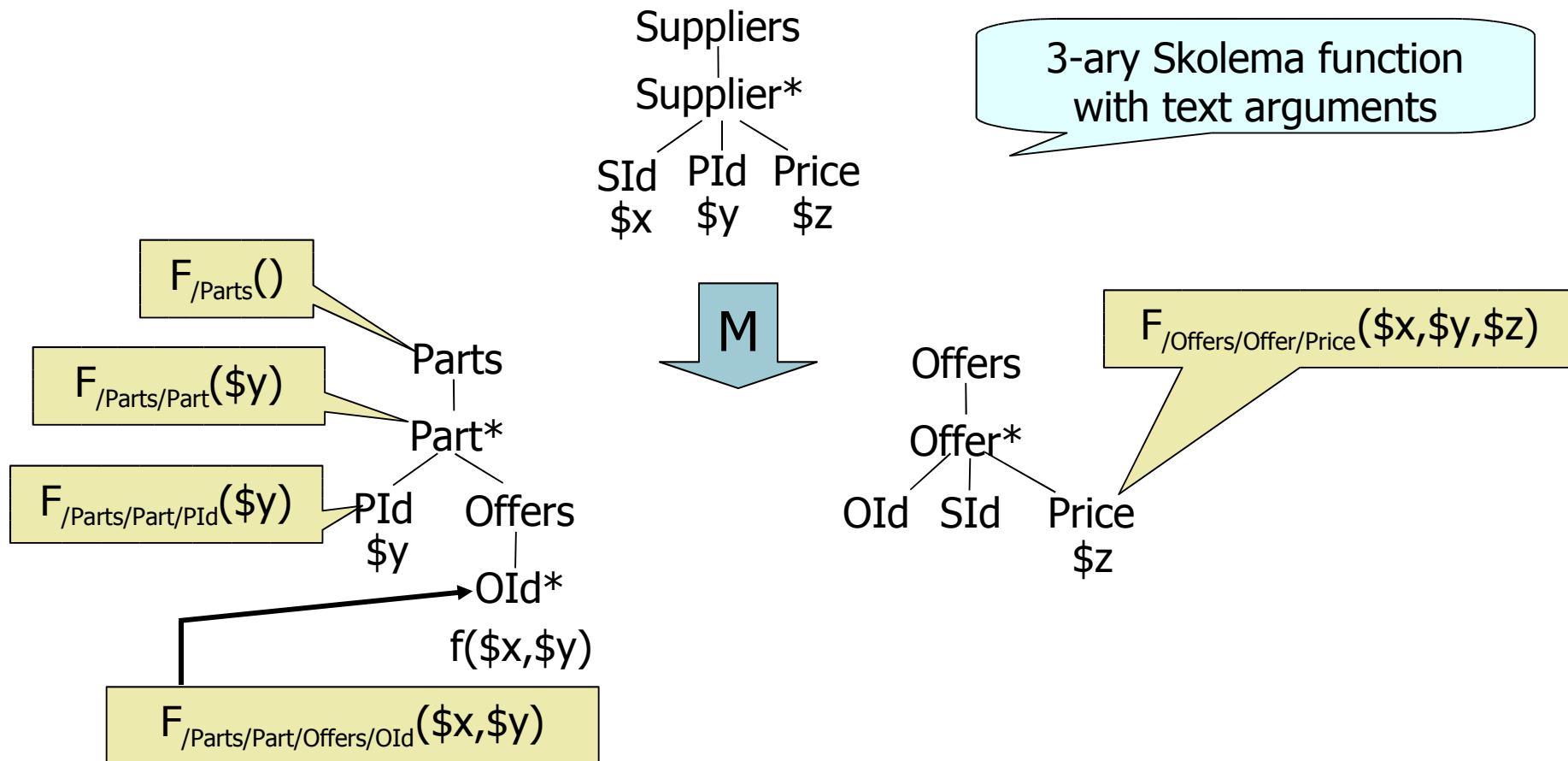


$$M \circ q = \tau_M(q)$$

Mappings

1. In both cases mappings between XML schemas must be established
2. What do we need to define mappings?
 - a language to express relationships between schemas (XDMap)
 - a method to create specification – as automatically as possible
 - mappings can be generated automatically from key constraints (within XML Schema)

The idea of mappings

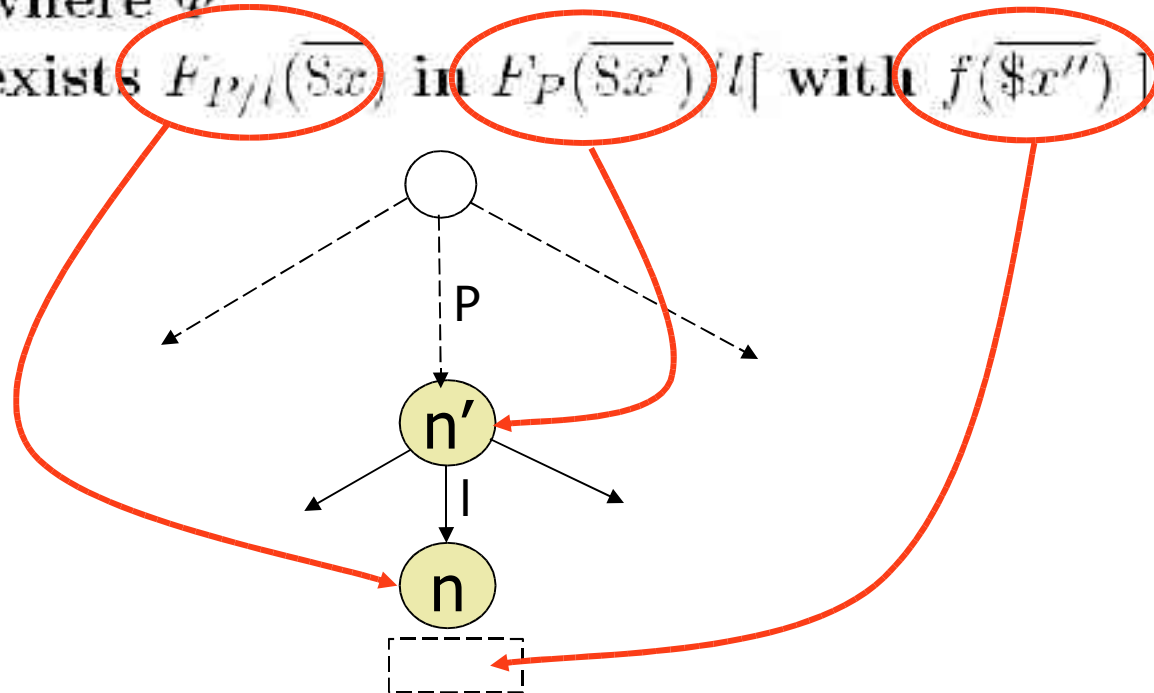


The idea of mapping specification

- $\forall \underline{x} (\Phi_S(\underline{x}) \Rightarrow \exists \underline{y} \Psi_T(\underline{x}, \underline{y}))$ – *data generating dependency*.
- $\forall \underline{x} (\Phi_S(\underline{x}) \Rightarrow \Psi_T(\underline{x}))$ – *full data generating dependency*.

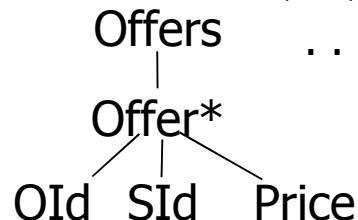
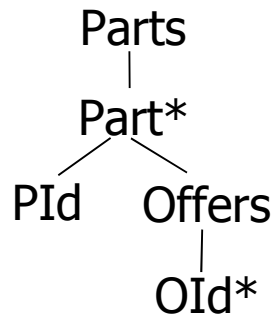
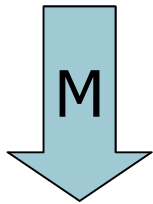
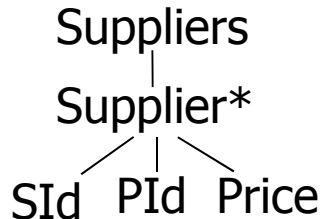
$M := \text{foreach } g, \dots, g$
 where Φ

exists $F_{P/l}(\overline{Sx})$ in $F_P(\overline{Sx'})/l$ with $f(\overline{Sx''})$



XDMap – a language for mapping specification

$$\forall \underline{x} (\Phi_S(\underline{x}) \Rightarrow \Psi_T(\underline{x}))$$



M = foreach \$s in /Suppliers/Supplier, \$x in \$s/SId,
\$y in \$s/PId, \$z in \$s/Price

where true
exists

$F_{/Parts/Part}(\$y)$ **in** $F_{/Parts}()/Part$

$F_{/Parts/Part/PId}(\$y)$ **in** $F_{/Parts/Part}(\$y)/PId$ **with** \$y

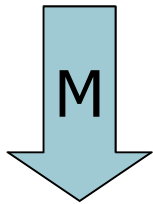
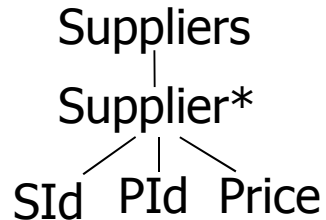
$F_{/Parts/Part/Offers}(\$y)$ **in** $F_{/Parts/Part}(\$y)/Offers$

$F_{/Parts/Part/Offers/OId}(\$x, \$y)$ **in** $F_{/Parts/Part/Offers}(\$y)/OId$ **with** $f(\$x, \$y)$

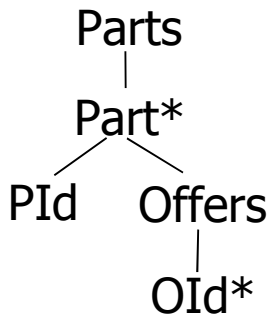
...

Sample target query (subset of XQuery)

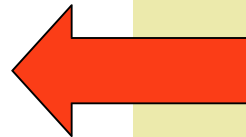
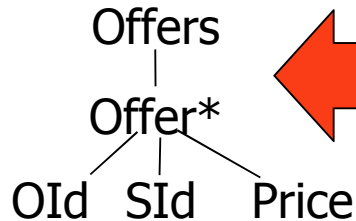
@S1:



@P3:



@O3:



T = (@P3, @O3), **S** = (@S1)

```
q =  
<Result>  
for $p in @P3/Parts/Part  
  $q in $p/Offers  
  $o in @O3/Offers/Offer  
where  
  $q/Old = $o/Old  
return  
  <Part>  
    <PartId> $p/PId </PartId>  
    <Supplier> $o/SId </Supplier>  
    <Price> $o/Price </Price>  
  </Part>  
</Result>
```

4-types of query components

```
q =  
<Result>  
for $p : @P3/Parts/Part }  
    $q : $p/Offers }  
    $o : @O3/Offers/Offer  
where  
    $q/Ord = $o/Ord }  
return  
    <Part>  
        <PartId> $p/PId </PartId>  
        <Supplier> $o/SId </Supplier>  
        <Price> $o/Price </Price>  
    </Part>  
</Result>
```

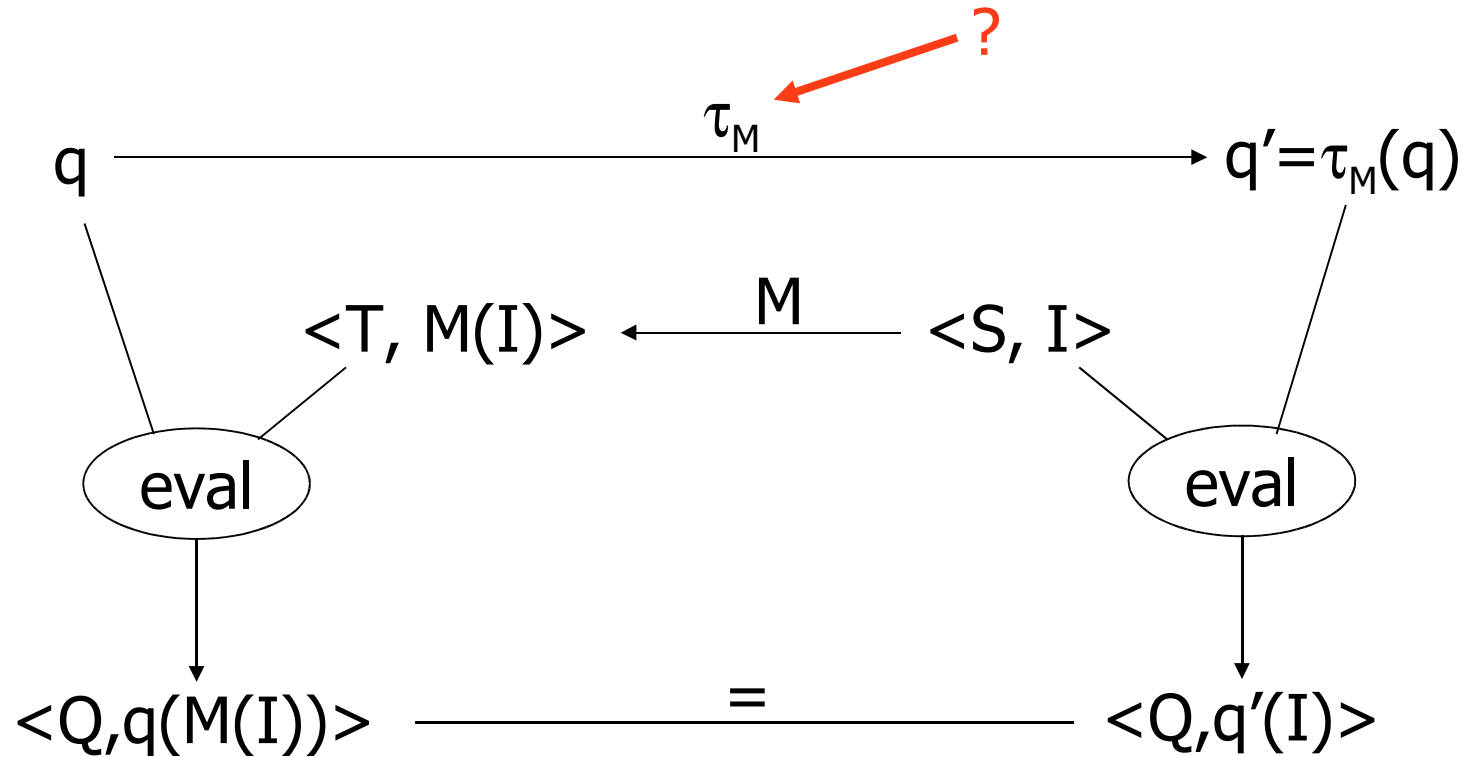
variable definition with absolute path

variable definition with relative path

comparison

text-valued expression

Query reformulation



Rewriting rule (R1-1)

```
q =  
<Result>  
for    $p in @P3/Parts/Part  
... 
```

$F_t : \$p \text{ in } @P3/Parts/Part$
$M^e : F_{@P3/Parts/Part}(\$y)$
$M^f : \$y \text{ in } @S1/Suppliers/Supplier/PId$
$\omega : [\$p \rightarrow \$p']$
$F_s : \$p' \text{ in } @S1/Suppliers/Supplier/PId$

```
q =  
<Result>  
for    $p' in @S1/Suppliers/Supplier/PId  
... 
```

Rewriting rule (R2-1)

```
q = <Result>
  for $p in @P3/Parts/Part
    $q in $p/Offers
    . . .
```

$F_t : \$q \text{ in } \p/Offers

$M^e : F_{@P3/Parts/Part/Offers}(\$s/PId) \text{ in } F_{@P3/Parts/Part}(\$s/PId)/\text{Offers}$

$M^f : \$s \text{ in } @S1/Suppliers/Supplier$

$\omega : [\$p \rightarrow \$p']$

$\omega : [\$q \rightarrow \$q']$

$F_s : \$q' \text{ in } @S1/Suppliers/Supplier$

$F_w : \$p' = \q'/PId

```
q = <Result>
  for $p' in @S1/Suppliers/Supplier/PId
    $q' in @S1/Suppliers/Supplier
    . . .
  where $p' = $q'/PId and ...
```

Rewriting rule (R1-2)

```
q = <Result>
  for $p in @P3/Parts/Part
    $q in $p/Offers
    $o in @O3/Offers/Offer
    . . .
```

```
Ft : $o in @O3/Offers/Offer
Me : F@O3/Offers/Offer($s/SId,$s/PId)
Mf : $s in @S1/Suppliers/Supplier
ω : [$o → $o' ]
Fs : $o' in @S1/Suppliers/Supplier
```

```
q = <Result>
  for $p' in @S1/Suppliers/Supplier/PId
    $q' in @S1/Suppliers/Supplier
    $o' in @S1/Suppliers/Supplier
    . . .
  where $p'=$q'/PId and ...
```

Process of query reformulation

We will say that a source query q_s is derived from a target query q_t by means of a reformulation τ_M , denoted by

$$(q_t, \emptyset) \xRightarrow[\tau_M]{*} (q_s, \Omega^*),$$

iff there is a sequence

$$(q^0, \Omega^0), \dots, (q^N, \Omega^N),$$

such that:

$$\begin{aligned} (q^0, \Omega^0) &= (q_t, \emptyset), \\ (q^i, \Omega^i) &\xRightarrow[\tau_M]{*} (q^{i+1}, \Omega^{i+1}) \\ (q^N, \Omega^N) &\equiv (q_s, \Omega^*), \\ &, 0 \leq i \leq N-1. \end{aligned}$$

Source query (reformulated q)

```
q' =  
<Result>  
for $p' : @S1/Suppliers/Supplier/PId  
    $q' : @S1/Suppliers/Supplier  
    $o' : @S1/Suppliers/Supplier  
where  
    $p' = $q'/PId and $q'/Old = $o'/Old and $q'/SId = $o'/SId  
return  
    <Part>  
        <PartId> $p' </PartId>  
        <Supplier> $o'/SId </Supplier>  
        <Price> $o'/Price </Price>  
    </Part>  
</Result>
```

Conclusion

1. E-commerce activities performed at electronic marketplace and using independently created e-commerce applications require data integration.
2. Data integration can be achieved by data transformation or query reformulation. In both cases mappings between schemas are necessary.
3. We propose a language XDMap for specifying mappings. Mappings are used for query reformulation.
4. Rules for query reformulation are proposed.

Thank You

