# Dynamic Model Harmonization between Unknown eBusiness Systems 

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## Agenda

* Introduction
* Concept of "Model Harmonization"
* Model Description
* Model Harmonization Algorithm
* Experiment
* Discussion on Business Resource Definition Matching
* Conclusions


## Introduction

## * Web Services

$\square$ from simple request-response to message-based conversation
$\square$ SOAP, WSDL, ebXML messaging, WS-Reliability, BTP
$\square$ Key issue: alignment of business process model among eBusiness systems

* Business process modeling technologies

■ WS-BPEL, WS-CDL, BPNM, BPML, ...
Usual assumption
$\square$ systems share an agreed business process model
■ not always be satisfied

* Approach in this research
$\square$ systems do not share complete business process models
- dynamically harmonize business process models
$\square$ to enable business conversation among independent and autonomous business systems using their best efforts


## * Usual assumption



- MDA
- CIM-PIM-PSM
- Repository
- Business process description lang.
- ...
- Model

Harmonization process model
$\qquad$

## * Two basic issues

- Alignment of business protocol


## main focus

■ Alignment of semantics of business resources in messages
separate issue


## Concept of "Model Harmonization"

* Focus on exposed models.
$\square$ models exposed by systems to other systems.
* An exposed model consists of an interface model and a behavior model.
* A type of business element in messages is specified by a $B R D$ (business resource definition).
■ product id, price, invoice, shipment notification, ...
$\square$ specified by a $B R D$ identifier, typically URI.
* Model Harmonization is a process to automatically adjust exposed models between encountering systems.


## Model Description

## * Formal definition of (exposed) models

```
M=(IM,BM)
    IM : An interface model defined below.
    BM : A behavior model defined below.
```

```
IM=(U,T,D)
    U}\mathrm{ : A finite set. (Operations)
    T:U->{I,O,(I,O),(O,I)}
    D:U\times{I,O }}->{(BRD,\ldots,BRD)
```

$B M=(B)$
$B$ : A regular expression on $U$.

■ Example of exposed model

■ Example mapping to WSDL style

```
U = {PriceQuery, GetPrice, Order, Response, Bill, PaymentNotice, Delivery, Acceptance }
T(PriceQuery)=O,T(GetPrice)=I, T(Order)=O,T(Response)=I,
T(Bill)=I, T(PaymentNotice)=O,T(Delivery)=I, T(Acceptance)=O
D(PriceQuery,O)= 'ns:price.ask', D(GetPrice,I)='ns:price.ans',
D(Order,O)='ns:form.order', T(Response,I)='ns:form.order_response',
D(Bill,I)='ns:form.bill', D(PaymentNotice,O)='ns:notice.conf',
D(Delivery,I)='ns:notice.shipment', D(Acceptance,O)='ns:notice.conf'
B = (SEQ,
        (OPT, (SEQ, PriceQuery, GetPrice) ),
        Order, Response,
        (OPT, ( SEQ, Bill, PaymentNotice, Delivery, Acceptance ) ) )
```

```
<definitions name= "Model_PS" xmlns:ns="http://www.examp.org/brd008/">
    <message name="MessageI">
        <part name="pt_1" element="ns:price.ask"/>
    </message>
    <message name="Message2">
        <part name="pt 1" element="ns:price.ans"/>
    </message>
    <portType name="Model_PS_interface">
        <operation name="P\overline{riceqQuery">}
            <output message="Message1">
        </operation>
        <operation name="GetPrice">
            <input message="Message2">
        </operation>
    </poritype>
    <behavior>
        <SEQ>
            <OPT>
            <SEQ> "PriceQuery" "GetPrice" </SEQ>
            </OPT>
            "Order" "Response"
            <OPT>
                <SEQ> "Bill" "PaymentNotice" "Delivery" "Acceptance" </SEQ>
                </OPT>
        </SEQ>
    </behavior>
</definitions>
```

■ Example of exposed model

■ Example mapping to WSDL style

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D(Order,O)='ns:form.order', T(Response,I)='ns:form.order_response',
D(Bill,I)='ns:form.bill', D(PaymentNotice,O)='ns:notice.conf',
D(Delivery,I)='ns:notice.shipment', D(Acceptance,O)='ns:notice.conf'
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    ( OPT, (SEQ, PriceQuery, GetPrice) ),
    Order, Response,
    ( OPT, ( SEQ, Bill, PaymentNotice, Delivery, Acceptance ) ) )
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        <part name="pt_1" element="ns:price.ask"/>
    </message>
    <message name="Message2">
        <part name="pt 1" element="ns:price.ans"/>
    </message>
    <pör\dot{tType name="Model PS_interface">}
        <operation name="Pric\overline{CQuery">}
            <output message="Message1">
        </operation>
        <operation name="GetPrice">
            <input message="Message2">
        </operation>
    </pority\dot{e}>>
    <behavior>
        <SEQ>
            <OPT>
                <SEQ> "PriceQuery" "GetPrice" </SEQ>
            </OPT>
            "Order" "Response"
            <OPT>
                <SEQ> "Bill" "PaymentNotice" "Delivery" "Acceptance" </SEQ>
            </OPT>
        </SEQ>
    </behavior>
</definitions>
```


## Model Harmonization Algorithm



* Top-level algorithm

Step a: exchange $P$ and $Q$ between $S$ and $T$
Step b: At $S$,
reduce $P$ adjusting to $Q$, using the model reduction algorithm suspend execution if the resulted model new $P$ is $\phi$
Step c: At $T$,
reduce $Q$ adjusting to $P$, using the model reduction algorithm suspend execution if the resulted model new $Q$ is $\phi$
Step d: start conversation between $S$ and $T$ using the reduced models

## Model Reduction Algorithm

* Notations in this explanation

■ Models:
$>P=\left(I M_{P}, B M_{P}\right), I M_{P}=\left(U, T_{P}, D_{P}\right), B M_{P}=\left(B_{P}\right)$
$>Q=\left(I M_{Q}, B M_{Q}\right), I M_{Q}=\left(V, T_{Q}, D_{Q}\right), B M_{Q}=\left(B_{Q}\right)$
$>T_{P}(u)$ and $T_{Q}(v)$ are either $I$ or $O, D_{P}\left(u, T_{P}(u)\right)$ and $D_{Q}\left(v, T_{Q}(v)\right)$ have one BRD.
■ Illustration:


$$
\begin{array}{ll}
u(\in U): & \text { an operation in } P . \\
t\left(=T_{P}(u)\right): & \text { its message exchange } \\
& \text { pattern, i.e., } I \text { or } O . \\
d\left(=D_{P}\left(u, T_{P}(u)\right):\right. & \text { its BRD. } \\
v(\in V): & \begin{array}{l}
\text { an operation in } Q . \\
r\left(=T_{Q}(v)\right):
\end{array} \\
& \begin{array}{l}
\text { its message exchange } \\
e\left(=D_{Q}\left(v, T_{Q}(v)\right):\right. \\
\text { pattern, i.e., } I \text { or } O .
\end{array} \\
\text { its BRD. }
\end{array}
$$

* Step1: Find corresponding sets of operations, $\Phi_{P}(u)$ and $\Phi_{Q}(v)$

$\square u$ and $v$ "correspond" when: strict version
$>t=O$ and $r=I$ and match $(d, e)=T r u e$, or $t=I$ and $r=O$ and match $(e, d)=$ True
$\square$ match () evaluates BRD matching:
$>$ match $\left(B R D_{\text {from }}, B R D_{t o}\right)= \begin{cases}\text { True } & \left(\text { (if } I\left(B R D_{\text {from }}\right) \subset I\left(B R D_{t o}\right) \text { is detected. }\right) \\ \text { False } & \left(\text { (if } I\left(B R D_{\text {from }}\right) \not \subset I\left(B R D_{t o}\right) \text { is detected. }\right)\end{cases}$ Unknown (otherwise, i.e., neither is detected.) Note: $I(b r d)$ denotes all instances satisfying brd.
$\square$ match () is implemented in various ways.
$>$ A trivial implementation:

$>$ More refined implementations
- will be discussed later.
* Step2: Remove unnecessary operations
$\square$ Remove operations not corresponding to any other operations, from behavior patterns $B_{P}$ and $B_{Q}$.
* Step3 (main part): Reduce behavior pattern
$\square$ Replace operations in a behavior pattern $B_{Q}$ by operations in $P$.
$\square$ for all $v$ in $B_{Q}$, replace $v$ by $G(v) . \quad\left(\right.$ call it $\left.B_{Q}^{\prime}\right)$


$$
G(v)=u_{1}\left|u_{2}\right| \ldots \mid u_{m}
$$

$\square$ Obtain $n e w B_{P}$ satisfying $L\left(n e w B_{P}\right)=L\left(B_{P}\right) \cap L\left(B_{Q}^{\prime}\right)$.
$>$ Note: $L(R)$ is a set of sequences matching to a regular expression $R$.
$\square$ Suspend if $L\left(n e w B_{P}\right)=\phi$.
Step4: Generate final model.

## Step 1:

for each $u$ in $U$ :

$$
\begin{aligned}
\Phi_{P}(u)=\{ & v \mid T_{P}(u)=O \text { and } T_{Q}(v)=I \text { and } \text { match }\left(D_{P}(u, O), D_{Q}(v, I)\right) \\
& \text { or } \left.T_{P}(u)=I \text { and } T_{Q}(v)=O \text { and } \text { match }\left(D_{Q}(v, O), D_{P}(u, I)\right)\right\}
\end{aligned}
$$

for each $v$ in $V$ :

$$
\Phi_{Q}(v)=\left\{u \mid T_{Q}(v)=O \text { and } T_{P}(u)=I \text { and match }\left(D_{Q}(v, O), D_{P}(u, I)\right)\right.
$$

$$
\text { or } \left.T_{Q}(v)=I \text { and } T_{P}(u)=O \text { and } \operatorname{match}\left(D_{P}(u, O), D_{Q}(v, I)\right)\right\}
$$

$U^{\prime}=\left\{u \mid u \in U\right.$ and $\left.\Phi_{p}(u) \neq \phi\right\} ; V^{\prime}=\left\{v \mid v \in V\right.$ and $\left.\Phi_{Q}(v) \neq \phi\right\}$
if $U^{\prime}=\phi$ or $V^{\prime}=\phi$ then exit (suspend)
Step 2:
obtain $B_{P}^{\prime}$ satisfying $L\left(B_{P}^{\prime}\right)=L\left(B_{P}\right) \cap\left(U^{\prime}\right)^{*}$ obtain $B_{Q}^{\prime}$ satisfying $L\left(B_{Q}^{\prime}\right)=L\left(B_{Q}\right) \cap\left(V^{\prime}\right)^{*}$ if $L\left(B_{f}^{\prime}\right)=\phi$ or $L\left(B_{Q}^{\prime}\right)=\phi$ then exit (suspend)

```
Step 3:
    for each v in V
        G(v)=\varepsilon
        for each u in }\mp@subsup{\Phi}{Q}{}(v):G(v)=G(v)+"|"+
    replace all v}\mathrm{ in }\mp@subsup{B}{Q}{\prime}\mathrm{ by }G(v
    obtain new }\mp@subsup{B}{P}{}\mathrm{ satisfying L (new }\mp@subsup{B}{P}{})=L(\mp@subsup{B}{P}{\prime})\capL(\mp@subsup{B}{Q}{\prime}
    if L (new\mp@subsup{B}{P}{})=\phi then exit (suspend)
```

Step 4:
$n e w U=\left\{u \mid u \in U^{\prime}\right.$ and $\left(u\right.$ in $\left.\left.n e w B_{P}\right)\right\}$
if $n e w U=\phi$ then exit (suspend)
new $T_{P}=T_{P}$ restricted in new $U$; new $D_{P}=D_{P}$ restricted in new $U$
$n e w I M_{P}=\left(n e w U, n e w T_{P}, n e w D_{P}\right) ; n e w B M_{P}=\left(n e w B_{P}\right)$
$n e w P=\left(\right.$ newIM $_{P}$, newBM $\left._{P}\right)$
return newP

## Experiment

* MD_REDUCE : an experimental implementation.
* Implements the model reduction algorithm.
* A model is represented by a Python object, independent of middleware platform. Mapping to WSDL style is supported.
* Execution example:
$\square$ A buyer's system exposing a model $P$ and a provider's system exposing a model $Q$.
$\square$ The trivial match () is used.


## Execution Example

## * Before reduction

## Model $P$

## NameSpaces:

ns1=http://www.examp.org/brd008/,
ns2=http://www.mkik.org/brdef/

## U, T, D:

PriceQuery[O]=ns2:price.ask, GetPrice[I]=ns2:price.ans, Order[O]=ns2:form.order,
Response[I]=ns2:form.order_response,
AnsQuestion[O]=ns2:form.order,
Bill[I]=ns1:form.bill, PaymentNotice[O]=ns1:notice.conf, Delivery[I]=ns1:notice.shipment,
Acceptance[O]=ns1:notice.conf
B:
(SEQ,
(OPT, (SEQ, PriceQuery, GetPrice)), Order,
(CHO,
Response, (SEQ, Response, AnsQuestion, Response), (SEQ,
(CHO,
Response,
(SEQ, Response, AnsQuestion, Response) ), Delivery, Acceptance, Bill, PaymentNotice )
) ) (

## Model $Q$

## NameSpaces:

NS=http://www.examp.org/brd008/, mk=http://www.mkik.org/brdef/

## U, T, D:

EstPrice[I]=mk:price.ask, AnsPrice[O]=mk:price.ans,
Preordered[I]=NS:notice.preorderID,
Order $[\mathrm{I}]=\mathrm{mk}$ :form.order,
Question[O]=mk:form.order_response,
AcceptOrder $[\mathrm{O}]=\mathrm{mk}:$ form.order_response,
DenyOrder[O]=mk:form.order_response,
Bill[O]=NS:form.bill, ConfPayment[I]=NS:notice.conf
Ship[O]=NS:notice.shipment, ConfDelivery[I]=NS:notice.conf,
B:
(SEQ,
(CHO, (SEQ, EstPrice, AnsPrice), Preordered),
Order,
(REP, (SEQ, Question, Order)),
(CHO,
DenyOrder,
(SEQ, AcceptOrder,
(CHO, ( SEQ, Bill, ConfPayment, Ship, ConfDelivery), ( SEQ, Ship, ConfDelivery, Bill, ConfPayment ) ) )

## * After reduction

## Model $P$

## NameSpaces:

ns1=http://www.examp.org/brd008/,
ns2=http://www.mkik.org/brdef/

## U, T, D:

PriceQuery[O]=ns2:price.ask, GetPrice[I]=ns2:price.ans, Order[O]=ns2:form.order,
Response[I]=ns2:form.order_response,
AnsQuestion[O]=ns2:form.order,
Bill[I]=ns1:form.bill, PaymentNotice[O]=ns1:notice.conf, Delivery[I]=ns1:notice.shipment,
Acceptance[O]=ns1:notice.conf
B:
(SEO,
PriceQuery, GetPrice, Order,
(CHO,
Response,
(SEQ, Response, AnsQuestion, Response), (SEQ,
(CHO,
Response,
(SEQ, Response, AnsQuestion, Response) ), Delivery, Acceptance, Bill, PaymentNotice )
) ) (

## Model $Q$

## NameSpaces:

NS=http://www.examp.org/brd008/,
mk=http://www.mkik.org/brdef/
U, T, D:
EstPrice[I]=mk:price.ask, AnsPrice[O]=mk:price.ans,
Order[I]=mk:form.order,
Question[O]=mk:form.order_response,
AcceptOrder $[\mathrm{O}]=\mathrm{mk}:$ form.order_response,
DenyOrder[O]=mk:form.order_response,
Bill[O]=NS:form.bill, ConfPayment[I]=NS:notice.conf
Ship[O]=NS:notice.shipment, ConfDelivery[I]=NS:notice.conf,
B:
(SEQ,
EstPrice, AnsPrice,
Order,
(CHO,
DenyOrder,
(SEQ, Question, Order, DenyOrder),
(SEQ,
(CHO,
AcceptOrder,
(SEQ, Question, Order, AcceptOrder) ),
Ship, ConfDelivery, Bill, ConfPayment )

## Simulation Results

* Intermediate protocol errors are decreased.

|  | without model reduction |  |  | with model reduction |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | System 1 |  | System 2 |  | System 1 |  | System 2 |  |
|  | total | error | total | error | total | error | total | error |
| Case 1 | 72 | 48 | 24 | 0 | 24 | 0 | 24 | 0 |
| Case 2 | 12 | 8 | 4 | 0 | 4 | 0 | 4 | 0 |
| Case 3 | 18 | 14 | 24 | 20 | 4 | 0 | 4 | 0 |
| Case 4 | 16 | 14 | 18 | 16 | 2 | 0 | 2 | 0 |
| Case 5 | 4 | 4 | 4 | 4 | 0 | 0 | 0 | 0 |

## Discussion on Business Resource Definition Matching

* Model reduction algorithm is independent of match().
$\square$ Even the trivial match() takes effect.
$\square$ More refined match() will bring further effect.
* Our initial proposal on improvement of match () :
$\square$ BRDs are defined in various forms, sometimes in natural language.
■ Introduce a BRD description as a BRD's proxy, using Web Ontology technology.
- A BRD description includes two parts.
$>$ Binary relations: relationship between BRDs such as equivalence, subsumption(inclusion), disjointness(exclusion)
$>$ Structure : structural characteristics of BRD
■ We proposed an initial algorithm for match() traversing Internet.

```
<owl:Class rdf:about="http://www.examp.org/brd008/form.bill">
    <owl:equivalentClass rdf:resource="http://www.examp.org/brd008/form.seikyusho'/>
    <rdfs:subClassOf rdf:resource="http://www.mkik.org/form.request"/>
    <owl:disjointWith rdf:resource="http://www.examp.org/brd008/notice.shipment"/>
```

Binary relations
[rdfs:subClassOf](rdfs:subClassOf)
[owl:Restriction](owl:Restriction)
<owl:onProperty rdf:resource="http://www.struex.com/form.name"/>

## * Outline of a proposed BRD Matching Algorithm

Step 1: search recursively BRDs referable from $B R D_{\text {from }}$ using binary relations sort BRDs into appropriate sets $B_{f e}, B_{f s}$, or $B_{f d}$.
Step 2: search recursively BRDs referable from $B R D_{t o}$ using binary relations sort BRDs into appropriate sets $B_{t e}, B_{t s}$, or $B_{t d}$.
Step 3: make a finat decision inspeeting these six sorted sets.
Note: $B_{f e}, B_{f s}, B_{f d}, B_{t e}, B_{t s}$, and $B_{t d}$ are sets of BRDs.

- $B_{f e}$ as equivalent with $B R D_{f f} B_{f s}$ as subsumed by $B R D_{f}$, and $B_{f d}$ as disjoint with $B R D_{f}$
- $B_{t e}$ as equivalent with $B R D_{t}, B_{t s}$ as subsumed by $B R D_{t}$, and $B_{t d}$ as disjoint with $B R D_{t}$


## Details of the algorithm

## Step 1:

$B_{f e}=\{f r o m\} ; B_{f s}=\phi ; B_{f d}=\phi$
$\operatorname{sort}\left(\right.$ from, $\left.B_{f e}, B_{f f}, B_{f d}\right)$
Step 2:
$B_{t e}=\{t o\} ; B_{t s}=\phi ; B_{t d}=\phi$ $\operatorname{sort}\left(t o, B_{t e}, B_{t s}, B_{t d}\right)$
Step 3:
if $B_{f e} \cap B_{t e} \neq \phi$ then return $T$
if $B_{f s} \cap B_{t e} \neq \phi$ then return $T$
if $B_{f e} \cap B_{t d} \neq \phi$ then return $F$
if $B_{f s} \cap B_{t d} \neq \phi$ then return $F$
if $B_{f d} \cap B_{t e} \neq \phi$ then return $F$
if $B_{f d} \cap B_{t s} \neq \phi$ then return $F$
if $B_{f e} \cap B_{t s} \neq \phi$ then return $U$
if $B_{f s} \cap B_{t s} \neq \phi$ then return $U$
if $B_{f d} \cap B_{t d} \neq \phi$ then return $U$
$\operatorname{sort}(b, E, S, D)$
if not visible( $b$ ) then return
for each $c \in$ equivalentClass $(b)$ and $c \notin E \cup S \cup D$ : if $b \in E$ then add $c$ to $E ; \operatorname{sort}(c, E, S, D)$ if $b \in S$ then add $c$ to $S ; \operatorname{sort}(c, E, S, D)$
if $b \in D$ then add $c$ to $D ; \operatorname{sort}(c, E, S, D)$
for each $c \in \operatorname{subClass} O f(b)$ and $c \notin E \cup S \cup D$ :
if $b \in E$ then add $c$ to $S ; \operatorname{sort}(c, E, S, D)$
if $b \in S$ then add $c$ to $S$; $\operatorname{sort}(c, E, S, D)$
for each $c \in \operatorname{disjointWith}(b)$ and $c \notin E \cup S \cup D$ : if $b \in E$ then add $c$ to $D ; \operatorname{sort}(c, E, S, D)$ if $b \in S$ then add $c$ to $D ; \operatorname{sort}(c, E, S, D)$
return
Note 1: visible (brd) is a function to resolve if the description $b r d$ is referable.
Note2: Parameters $E, S$ and $D$ are used as "call by reference".

## Conclusions

* The concept of "Model Harmonization" was proposed.
* The formal definition of an exposed model was given.
* The algorithm to harmonize exposed models was provided.
* Its appropriateness and effectiveness were proved through experimental implementation.
* BRD matching was discussed and initial solution was proposed.
* Future works include:
$\square$ expansion of behavior pattern expression.
$\square$ internal error handling
$\square$ improvement of match().


## END

## Thank you very much.

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